## 6.5896 - Algorithmic Statistics

Lecture 9: The role of Temperature (part L)

Kesten-Stigum Bound

Setup: Broadcasting Model on a Tree

- complete binary tree Tof depth h

to other degrees a irregular trees

root node p

sedges directed away from p

- Random Process: samples r.v. Xv = {±1} at each v = V

Xp war {±1}
 processing all other nodes v following a topological ordering of T, sample Xv conditioning on Xu as follows

 $Pr[X_{v}=s \mid X_{u}=s]=1-p, \forall s \in \{\pm 1\}$   $aka Pr[X_{v}=-s \mid X_{u}=s]=p$ 

USSUME PE[0, 1] makes exposition slightly easier

(be is w.l.o.g. up to flipping 100-1
every offer layer)

Lythus each edge: (1-P P)

Channel M= (1-P P)

Lythings Ly things extend to - Very convenient decomposition of M  $\begin{pmatrix} 1-p & p \\ p & 1-p \end{pmatrix} = \begin{pmatrix} 1-2p \end{pmatrix} \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} + 2p \begin{pmatrix} 1/2 & 1/2 \\ 1/2 & 1/2 \end{pmatrix} -$ Lo se cond largest eigenvalue of M
corresponding to eigenvector ("1/5")
(the other eigenvalue is 1 under this event ne'll say edge "doesn't / corresponding to (1/5) 

with prub. 1-2=2p: X, w [±1]

under this event independently

we'll say edge' breaks" at X.

Observation: Suppose that X- E [±1] is a sample

from broadcast process on tree

then  $X_{V} \sim P_{\theta}(x) \propto exp(\frac{5}{cu,v) \in E}$ where  $\theta = \frac{1}{2} log \frac{1}{P}$  tree structured Ising model w/ same interaction strength

on all edges it no external fields

$$\left[ \begin{array}{c} \rho(X_{v} = \lambda_{v}) = \frac{1}{2} \cdot \prod_{(u,v) \in E} (1-p)^{\frac{1}{2} \lambda_{u} = \lambda_{v}} \cdot p^{\frac{1}{2} \lambda_{u} \neq \lambda_{v}} \\
= \frac{1}{2} \cdot \prod_{(u,v) \in E} (1-p)^{\frac{1+\lambda_{u} \times v}{2}} \cdot p^{\frac{1-\lambda_{u} \times v}{2}} \\
= \frac{1}{2} \cdot \prod_{(u,v) \in E} \ell \times \rho \left( \frac{1}{2} \left( 1 + \lambda_{u} \times v \right) \log (1-p) + \frac{1}{2} \left( 1 - \lambda_{u} \times v \right) \log p \right) \\
= \lambda \cdot \prod_{(u,v) \in E} \ell \times \rho \left( \frac{1}{2} \times u \times v \log \frac{1-p}{p} \right)$$

Things to note: 
$$P = \frac{1}{2} \implies \theta = 0$$
 all nodes independent
$$P = 0 \implies \theta = 1 \text{ and index indivitely}$$

p = 0 => 0 = +00 all nodes infinitely correlated

$$p = 0 \Rightarrow \theta = +\infty$$
 all nodes indinitely correlated
$$\begin{array}{c}
Correlated \\
Cir = 1/2
\end{array}$$

$$\begin{array}{c}
X : y = 1/2
\end{array}$$

~> | L(K) | = 2h

· Question: Suppose L(h) are nodes at depth h.

How much information about xo does Xully contain?

Observation 1: with probability  $(2p)^2$ , children of the root & therefore all their descendents are independent of root!

If this happens already at depth 1, what happens as h-> =>?

To study this formally:

Suppose Mt. (·) = Pr[XL(R)=· | Xp=+1]

& 4- (.) = Pr[X1(A)=. | Xp=-1]

Le Cam => given XL(R) the best estimator of the root state has error

$$\frac{1}{2}\left(1-TV(\mu_{R}^{+},\mu_{\bar{A}}^{-})\right)$$
Define the black production could

Def: We'll say that the root reconstruction problem is solvable if

 $\lim_{h\to\infty}\inf TV(\mu_h^+,\mu_h^-)>0$ o.w. we call it unsolvable

Observation 2: Reconstruction problem is unsolvable if 2.2<1

Lexpected # of Children who copy

The Children who copy

The Lexpected # of Children who copy

The Lexpected # of Children who copy

The Children who copy

The

By birth-heath process analysis if 21<1 the state of the root will "die" w pr 1 on un infinite tree

nodes that copy Xo

nodes that are independent of Xo

Question: Is it solvable when  $2\lambda > 1 \iff p < \frac{1}{4}$ ?

Theorem. The root reconstruction problem is so vable iff  $2\lambda^2 > 1$  ( $\Leftrightarrow p < \frac{1}{4}(2-\overline{l2})$ )

0.5858 Proof: we will only show sufficiency
i.e.  $2\lambda^2 > 1 \implies$  solvable

· We'll show that the majority of the leaf values maintains reconstruction error < 2 as h >>>

OLemma 1: E[Ze | Xe] = Xe  $\mathbb{E}\left[Z_{h} \mid \chi_{g}\right] = \frac{1}{1^{h}} \mathbb{E}\left[X_{l} \mid \chi_{g}\right]$   $\mathbb{E}\left[X_{l} \mid \chi_{g}\right]$   $\mathbb{E}\left[X_{l} \mid \chi_{g}\right]$  $=\frac{1}{\lambda^{h}}\cdot\left(\chi_{g}\cdot\lambda^{h}+0\cdot\left(1-\lambda^{h}\right)\right)=\chi_{g}$ It any of these edges "break" Xe, is independent If home of the edges "break" Xe = Xo  $\frac{1/2}{1-(2l^2)^{-1}}, \text{ if } 2l^2 > 1$   $\frac{1}{1-(2l^2)^{-1}}, \text{ if } 2l^2 > 1$ , if 22° < 1 Proof: postpone for a bit mental picture:

Using Lemma 1 & Lemma 2 to prove sufficiency when  $2\lambda^2 > 1$ Te call definitions of  $\mu_h^+$ ,  $\mu_h^-$ ,  $\mu_h^+$ ,  $\mu_h^-$ Claim:  $TV(\hat{\mu}_h^+, \hat{\mu}_h^-) \in TV(\hat{\mu}_h^+, \hat{\mu}_h^-)$ proof: easy  $Z_h = f(X_{L_h})$ Caverage f'n

take optimal coupling of  $X_{L_h}^{\dagger} \sim \mu_h^{\dagger}$  and  $X_{L_h} \sim \mu_h^{\dagger}$ 

implies a coupling of  $Z_h^+ \sim \hat{\mu}_h^+$  and  $Z_h^- \sim \hat{\mu}_h^-$ 

under this coupling:  $\Re \left[ Z_h^{\dagger} \neq Z_h^{\dagger} \right] \leq \Re \left[ X_{L_h}^{\dagger} \neq X_{L_h}^{\dagger} \right]$   $\left[ V(\hat{y}_h^{\dagger}, \hat{y}_h^{\dagger}) \right] = 0$ To  $\left[ V(\hat{y}_h^{\dagger}, \hat{y}_h^{\dagger}) \right] = 0$ 

TV(ph, ph)

15 some

Coupling

Of ph and coupling of

Ph Ph, ph

To show solvability, suffices to lower bound

$$=\frac{1}{4}\frac{\left(\mathbb{E}_{h}^{+}(z_{h})-\mathbb{E}_{h}^{-}(z_{h})\right)^{2}}{Var(z_{h})}$$

$$=\frac{1}{Var(z_{h})}$$

$$=\frac{1}{Var(z_{h})}$$

$$=\frac{1}{Var(z_{h})}$$

$$\Rightarrow\lim_{h\to\infty}\inf Tv(\mu_{h}^{+},\mu_{h}^{-})$$

$$\Rightarrow\lim_{h\to\infty}\inf Tv(\mu_{h}^{+},\mu_{h}^{-})$$

$$\Rightarrow\lim_{h\to\infty}\inf Var(z_{h})$$

$$\Rightarrow\lim_{h\to\infty}\frac{1}{Var(z_{h})}$$

 $Var[Z_h] = Var[\{[Z_h|X_g]\}] + \{[Var[Z_h|X_g]\}$ 

=  $Var\left[\chi_{g}\right] + \frac{1}{2} Var\left[\chi_{g}=1\right]$ 

+ 1 Var[ = 1 | X = -1]

 $\left(\begin{array}{c} z \\ \overline{z} \\ \overline{z$ 

2 2° /h/(2)

using 
$$Var[Xg]=1$$

be symmetry

To understand  $Var[Z_h|X_p=+1]$ 

To understand  $Var[Z_h|X_p=+1]$  consider

$$Z_h = Z_h^{(1)} + Z_h^{(2)}$$

 $Z_{h} = Z_{h}^{(1)} + Z_{h}^{(2)}$   $\Rightarrow Contribution of leaves of T_{1}$   $\Rightarrow Contribution of leaves of T_{2}$   $\Rightarrow Contribution of leaves of T_{2}$   $\Rightarrow Contribution of leaves of T_{3}$   $\Rightarrow Contribution of leaves of T_{4}$   $\Rightarrow Contribution of leaves of T_{4}$ 

$$Var\left[\frac{2}{2}h\left|\chi_{p}=+1\right] = Var\left[\frac{2}{2}h^{(1)}\left|\chi_{p}=+1\right]+Var\left[\frac{2}{2}h^{(2)}\left|\chi_{p}=+1\right]\right]$$

$$= 2 Var\left[\frac{2}{2}h^{(1)}\left|\chi_{p}=+1\right]$$

Symmetry = 
$$2 \cdot \left( \mathbb{E} \left[ \left( \mathbb{Z}_{h}^{(i)} \right)^{2} | \chi_{p}=+1 \right] - \left( \mathbb{E} \left[ \mathbb{Z}_{h}^{(i)} | \chi_{p}=+1 \right]^{2} \right)$$

Now note:  $\mathbb{E} \left[ \mathbb{Z}_{h}^{(i)} | \chi_{p}=+1 \right] = \frac{1}{2}$  (by symmetry b.w. they have the same

have the same expectation cond. on 70=+1 & the sum of these expectation

is +1)

 $\mathbb{E}\left(\left(\mathcal{Z}_{h}^{(1)}\right)^{2} \middle| \mathcal{X}_{p}=+1\right)=\left(1-p\right) \cdot \mathbb{E}\left(\left(\mathcal{Z}_{h}^{(1)}\right)^{2} \middle| \mathcal{X}_{g^{(1)}}=+1\right)$ + p. [ (2h) 2 | Xp(1) =-1]

= 
$$(1-p) \cdot \frac{1}{(21)^2} \mathbb{E} \left[ Z_{h-1}^2 | \chi_p = +1 \right] + p \cdot \frac{1}{(21)^2} \cdot \mathbb{E} \left[ Z_{h-1}^2 | \chi_p = -1 \right]$$

$$\mathbb{E} \left[ Z_{h-1}^2 | \chi_p = +1 \right]$$

$$= \frac{1}{(21)^2} \cdot \mathbb{E} \left[ Z_{h-1}^2 | \chi_p = +1 \right]$$

putting everything together:

$$Var[Z_{h}] = 1 + 2 \cdot \left(\frac{1}{(2\lambda)^{2}} \mathbb{E}[Z_{h-1}^{2} | \chi_{p}=+1] - \frac{1}{4}\right)$$

$$= 1 + \frac{1}{2\lambda^{2}} \cdot \mathbb{E}[Z_{h-1}^{2} | \chi_{p}=+1] - \frac{1}{2}$$

$$= \frac{1}{2} + \frac{1}{2\lambda^{2}} \cdot \left( Var[Z_{h-1} | \chi_{p}=+1] + \left( \frac{1}{2} \left[ Z_{h-1} | \chi_{p}=+1 \right]^{2} \right) \right)$$

$$= \frac{1}{2} + \frac{1}{2\lambda^{2}} \cdot \left( Var[Z_{h-1} | \chi_{p}=+1] + 1 \right)$$

$$= \frac{1}{2} + \frac{1}{212} \cdot \left( \text{Var}[Z_{h-1}] \times p = +1 \right) + 1$$

$$= \frac{1}{2} + \frac{1}{212} \cdot \text{Var}[Z_{h-1}] \text{ as shown}$$

$$= \frac{1}{2} + \frac{1}{212} \cdot \text{Var}[Z_{h-1}]$$

Solving recursion gives result &

$$= a + b \left( a + b g(h-2) \right)$$

$$= a + b \cdot a + b^{2} \cdot g(h-2)$$

$$= a + b \cdot a + b^{2} \cdot a + b^{3} \cdot g(h-2)$$

$$= a \left( 1 + b + b^{2} + \dots + b^{k-1} \right) + b^{k} \cdot g(h-k)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a \left( 1 + b + \dots + b^{k-1} \right) + b^{k} \cdot g(h)$$

$$= a$$

g(h) = a + b · g(h-1)

If 
$$2h^2=1$$
 $g(h)=a+g(h-1)=h\cdot a+g(0)$ 

$$Var[Z_h] = h \cdot \frac{1}{2} + 1 \rightarrow t\infty$$

Post Morton:

high temperature 
$$2\lambda^2 < 1$$
  $\Leftrightarrow$   $p > \frac{1}{z}(1 - \frac{1}{12})$ 

$$0 < \frac{1}{z} \log \frac{\sqrt{z+1}}{\sqrt{z-1}}$$
is BAD: prexents root reconstruction

low temperature 
$$2\lambda^2 > 1 \Leftrightarrow p < \frac{1}{2}(1-\frac{1}{12}) \Leftrightarrow p > \frac{1}{2}\log\frac{[2+1]}{[2-1]}$$

is GOOD: allows root reconstruction

lim sup  $\Pr$  [incorrect reconstruction] =  $h \to \infty$ 

lim sup  $\frac{1}{2}(1-TV(\mu_{h+1}^+\mu_h^-)) = \frac{1}{2}-(1-\frac{1}{2\lambda^2}) < \frac{1}{2}$ 
 $h \to \infty$ 

- o d-ary tree, binary symmetric channel: threshold at  $5.\lambda^2 = 1$
- more general setting: q possible states per mode d-ary tree
   every edge: Channel M

• for q=2: threshold at  $d_2(m)^2=1$ d·l<sub>z</sub>(M)<sup>2</sup>>1 sufficient
[Kesten-Stigum<sup>1</sup>66] necessary [Bleher- Zuiz-Zagrebnov '95 ∂ larger q: d. 22(m)²>1 sufficient [Kesten-Stigum'66] but not necessary! [Mossel '01] [Sly'09] give examples for 9),5

for which  $\lambda_2(m)=0$  yet

reconstruction is possible! for q=3,4: KS-bound is tight

if d>do(m) [mossel-sly-sohn'23] o for any q, if reconstruction algo only uses
the counts of how many leaves are +1 or -1
then

KS-Loynd is tight! [Mossel-Pares'03] o Conjecture [Koehler-Mosse| '22]: not just linear fn's but
also low-degree polynomials of the

XLL also fail below K-S bound!